## PtCut: A Program To Compute Tropical Prevarieties

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- Input: a polynomial system of equations.
- (This system is derived from ODEs, which are derived from a (bio)chemical reaction network with mass-action kinetics.)
- Rate constants are symbolic, but values have to be specified for computation.
- Output: the tropical prevariety of the system.
- We want to deduce approximate knowledge about the classical variety from the study of its tropical prevariety.

- First, the system is tropicalized. This means, each '·' becomes a '+' and each '+' becomes a 'min'. There is no translation for '-'!
- The multiplicative constant (i.e., the replaced rate constants) is transformed through a log<sub>ε</sub>(x), ε < 1, and rounded to integer.</li>
- Example: the classical equation x<sub>1</sub> + 3x<sub>2</sub> + x<sub>3</sub><sup>5</sup> becomes min(x̂<sub>1</sub>, log<sub>ε</sub> 3 + x̂<sub>2</sub>, 5x̂<sub>3</sub>) in the tropical world.
- (By abuse of language, I'll often keep the variable names:  $Trop(x_4 + x_5^4) = min(x_4, 4x_5).)$
- Choice of  $\varepsilon$  and rounding precision has influence on the number and shape of solutions and run-time.

# Tropicalization (2)

- A classical root is found if f(x) = 0. That is, positive and negative monomials cancel each other out:  $\sum_{i} P_i = \sum_{j} N_j$ ,  $P_i$ ,  $N_j \ge 0$ .
- Since there is no tropical subtraction, a *tropical root* is defined as the points where the minimum of tropical monomials is attained at least twice:

$$P_k \ge P_i = N_j \le N_\ell. \tag{1}$$

- Cycling over all pairs of monomials and applying (1) defines a *tropical hypersurface*, which in turn defines a set of polyhedra.
- The intersection of the tropical hypersurfaces of multiple polynomials  $f_i$  is called a *tropical prevariety*.
- For "complex" solutions, we ignore the sign of monomials:

$$P_i = P_j \leq P_k, \ i \neq j.$$

- Polynomial systems often come from the Biomodels database (in which case a SBML-parser is needed to build the ODEs from SBML.) The database specifies rate constants as well.
- PtCut reads this polynomial system as a list of polynomials  $f_r$  in  $x_i$  with parameters  $k_j$ . PtCut tries to solve the system of equations  $f_r = 0$ .
- Parameters are replaced, then polynomials are tropicalized and thus create sets of polyhedra.
- These sets of polyhedra are intersected one by one to get the tropical prevariety.
- Some tricks help to keep the run-time low (we'll come to that).

This is the MAPK model (Markevich et al., 2004) as modeled in the BioModels database.



## Example: Biomodel 26 (2)

From the above network one can derive the following ODEs:

$$\begin{split} \dot{x}_1 &= k_2 x_6 + k_{15} x_{11} - k_{1x} x_4 - k_{16} x_{1x5} \\ \dot{x}_2 &= k_3 x_6 + k_5 x_7 + k_{10} x_9 + k_{13} x_{10} - x_2 x_5 (k_{11} + k_{12}) - k_4 x_2 x_4 \\ \dot{x}_3 &= k_6 x_7 + k_8 x_8 - k_7 x_3 x_5 \\ \dot{x}_4 &= x_6 (k_2 + k_3) + x_7 (k_5 + k_6) - k_1 x_1 x_4 - k_4 x_2 x_4 \\ \dot{x}_5 &= k_8 x_8 + k_{10} x_9 + k_{13} x_{10} + k_{15} x_{11} - x_2 x_5 (k_{11} + k_{12}) - k_7 x_3 x_5 - k_{16} x_1 x_5 \\ \dot{x}_6 &= k_1 x_1 x_4 - x_6 (k_2 + k_3) \\ \dot{x}_7 &= k_4 x_2 x_4 - x_7 (k_5 + k_6) \\ \dot{x}_8 &= k_7 x_3 x_5 - x_8 (k_8 + k_9) \\ \dot{x}_9 &= k_9 x_8 - k_{10} x_9 + k_{11} x_2 x_5 \\ \dot{x}_{10} &= k_{12} x_2 x_5 - x_{10} (k_{13} + k_{14}) \\ \dot{x}_{11} &= k_{14} x_{10} - k_{15} x_{11} + k_{16} x_{15} \end{split}$$

Some additional calculations yield the following conservation laws:

$$k_{17} = x_5 + x_8 + x_9 + x_{10} + x_{11}$$
  

$$k_{18} = x_4 + x_6 + x_7$$
  

$$k_{19} = x_1 + x_2 + x_3 + x_6 + x_7 + x_8 + x_9 + x_{10} + x_{11}$$

Setting all differentials to zero, we only keep the polynomials as input for PtCut. The parameters are just extracted from BioModels database.

Polynomial\_system.txt: Params txt  $k^{2}x^{6} + k^{1}5x^{11} - k^{1}x^{1}x^{4} - k^{1}6x^{1}x^{5}$ k1 = 0.02k3\*x6 + k5\*x7 + k10\*x9 + k13\*x10 - x2\*x5\*(k11 + k12) - k4\*x2\*x4 $k^2 = 1$ k6\*x7 + k8\*x8 - k7\*x3\*x5k3 = 0.01x6\*(k2 + k3) + x7\*(k5 + k6) - k1\*x1\*x4 - k4\*x2\*x4k4 = 0.032k8\*x8 + k10\*x9 + k13\*x10 + k15\*x11 - x2\*x5\*(k11 + k12) - k7\*x3\*x5 - k16\*x1\*x5k5 = 1 $k_{1*x_{1}x_{4}} - x_{6*}(k_{2} + k_{3})$ k6 = 15k4\*x2\*x4 - x7\*(k5 + k6)k7 = 0.045k7\*x3\*x5 - x8\*(k8 + k9)k8 = 1k9\*x8 - k10\*x9 + k11\*x2\*x5k9 = 0.092k12\*x2\*x5 - x10\*(k13 + k14)k10 = 1k14\*x10 - k15\*x11 + k16\*x1\*x5 k11 = 0.01x5 - k17 + x8 + x9 + x10 + x11k12 = 0.01x4 - k18 + x6 + x7k13 = 1x1 - k19 + x2 + x3 + x6 + x7 + x8 + x9 + x10 + x11k14 = 0.5 $k_{15} = 0.086$ k16 = 0.0011k17 = 100

> k18 = 50k19 = 500

## Combining All Polytopes (1)

• Each tropical polynomial gives rise to a tropical hypersurface that describes a set (a disjunction) of polyhedra:

$$B_i = \bigcup_j P_{ij}, \quad i \in [1:n].$$

• The intersection (conjunction) of all tropical hypersurfaces is the tropical prevariety *U*:

$$U=\bigcap_{i}^{n}B_{i}=\bigcap_{i}^{n}\bigcup_{j}P_{ij}.$$

• Because of distributivity, we can rewrite that as:

$$U = \bigcup_{\ell \in \mathcal{I}} \bigcap_{i}^{n} P_{i\ell_i}, \text{ with } \mathcal{I} = \bigotimes_{i}^{n} |B_i|.$$

- The number of *possible* combinations is exponential in *n*.
- However, the biological models we examined so far have usually a very small (< 10) number of polyhedra in the prevariety, with only one exception (Biomodel 102), where it is  $\approx$  400.
- Still, the number of *intermediate* polyhedra can be much higher. Keeping this number low is important for fast computations. This can be done by several methods:
  - If, after intersecting two hypersurfaces, for two of the resultant polyhedra  $P \subseteq Q$  holds, we keep only Q.
  - The order of evaluation is important! A good heuristic is to intersect hypersurfaces with the fewest polyhedra first.
  - If we find constraints that are shared by all polyhedra of a hypersurface, they must be shared by the polyhedra of the prevariety as well. Hence, we can apply those constraints to all unprocessed hypersurfaces to reduce the number of polyhedra in them.

## Time vs. Dimension



## Time vs. Possible Combinations



## PtCut vs. Satya's Solution



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PtCut: Compute Tropical Prevarieties

Grid Sampling:

- PtCut can run calculations over a grid of parameters values. Only changed equations are re-calculated to save time.
- The grid can be additional or multiplicative, i.e., linear or factorial. This is set independent for each parameter.
- The complete  $30 \times 18$  grid from (Bradford et al., 2017) takes 260 sec to compute, i.e., 0.4 sec per grid point.

Other:

 PtCut can calculate all solutions ("complex") or only the ones where the minimum is attained by monomials with opposite signs ("real"). The latter is much faster and yields fewer solutions.

• Specify 
$$\varepsilon = \frac{1}{X}$$
 with  $-eX$  and scaling by  $10^X$  with  $-rX$ .

- PtCut is written in Python 2.7 for SageMath, a free computer algebra system.
- Tests were run on a virtual machine under Ubuntu 16.04 TLS, 64-bit, with 24 GiB of memory.
- The CPU was an Intel i7-3930K with 3.20 GHz. The code does not actively make use of multithreading.
- PtCut can be started via terminal or SSH and produces only text mode output and log/solution files.
- PtCut source code is about 3000 lines of code and licensed under LGPL v3.
- Python is easy to read and well supported by libraries, so PtCut is easy to modify.

- What can we infer from the tropical prevariety?
- Make PtCut run independently from SageMath and possibly directly under Window, so it's easier to use.
- Complete an open implementation of an SBML-parser.
- Can we further improve the speed? E.g., Biomodel 19 with dimension 92 and  $\approx 10^{60}$  combinations is still out of reach.
- Add support for variable substitution with the vertex cover method.
- Turn a parameter into a variable to find its breaking point without grid searching.

PtCut web page: http://wrogn.com/ptcut

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